

Optimal Play of the Great Rolled Ones Game

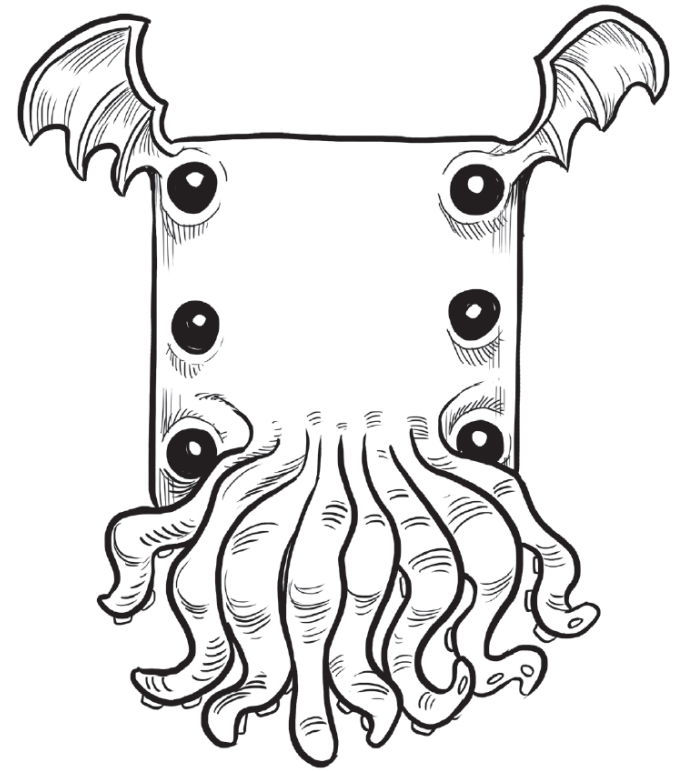
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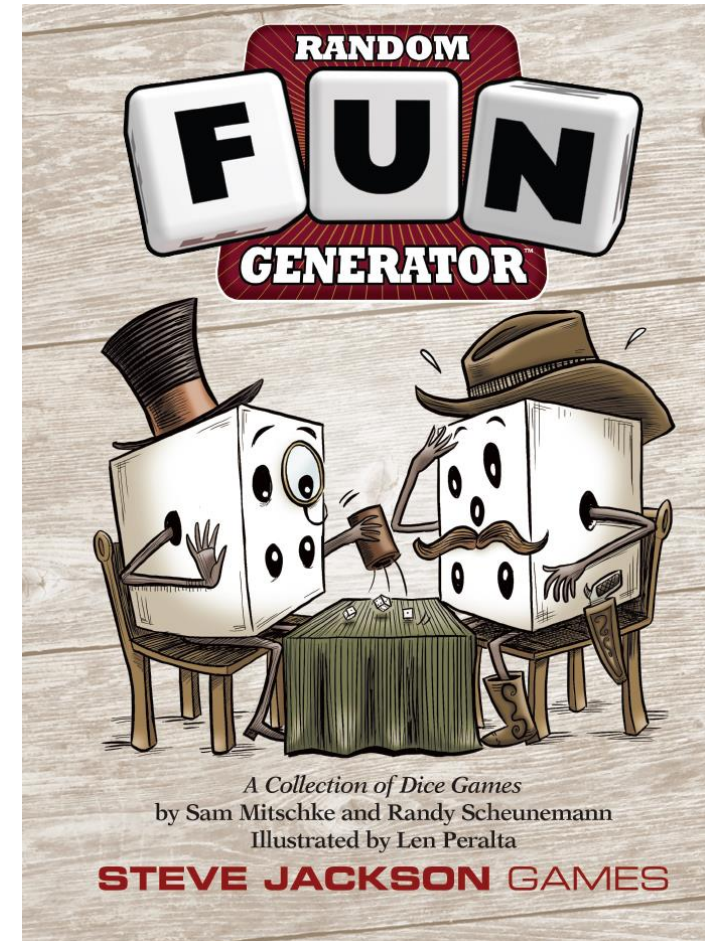
Overview

- Great Rolled Ones Rules
- Optimality Equations
- Solution Method
- Optimal Compensation Points and Visualization
- Human-Playable Policy: Simple player and ones cases
- Conclusions



Great Rolled Ones

- A **jeopardy dice game** for 2 or more players. (Here we consider 2 player only.)
 - Jeopardy (“Push your luck”) Dice Game – Primary mechanic: Roll/hold decisions where holding *secures* turn progress, whereas rolling *risks* all turn progress for potentially greater turn progress.
- First published in 2020 by Sam Mitschke and Randy Scheunemann
 - Similar to the dice game Zombie Dice
 - Both are jeopardy dice games in the Ten Thousand dice game family



Great Rolled Ones Rules



- 2 or more players using **5 standard (d6) dice**.
- Players will have the **same number of turns**. A turn consists of a **sequence of player dice rolls where rolled 1s are set aside**.
- The **turn ends when** either the player
 - decides to **hold** (i.e. stop rolling) and score the **total number of non-1s rolled**, or
 - has **rolled three or more 1s**, ending the turn with no score change.
- A round consists of each player taking one turn in sequence.
- Any player ending their turn with a **goal score of 50** or more causes that to be the last round of the game.
- At the end of the last round, the player with the highest score wins.
- (We assume that a player is constrained to attempt to exceed the score of the current leader in the last round.)

Great Rolled Ones Example Round

Player	Roll	Result (Decision)
1	1, 1, 3, 4, 5	Two 1s set aside, turn total 3 (roll)
1	2, 2, 4	No 1s set aside, turn total 6 (roll)
1	1, 1, 6	Two 1s set aside for a total of four 1s, \geq three 1s → turn ends with no score gain
2	4, 4, 4, 4, 5	No 1s set aside, turn total 5 (roll)
2	4, 4, 4, 5, 5	No 1s set aside, turn total 10 (roll)
2	1, 1, 2, 4, 5	Two 1s set aside, turn total 13 (hold) → turn ends with a score gain of 13

Optimality Equations: Probability of Rolling 1s

Nonterminal states are described as the 5-tuple (p, i, j, k, o) , where p is the current player number (1 or 2), i is the current player score, j is the opponent score, k is the turn total, and o is the number of rolled 1s set aside.

Let $P_{\text{new1s}}(d, o_{\text{new}})$ denote the probability that o_{new} of d dice rolled are 1s ($0 \leq o_{\text{new}} \leq d \leq 5$):

$$P_{\text{new1s}}(d, o_{\text{new}}) = \binom{d}{o_{\text{new}}} \left(\frac{1}{6}\right)^{o_{\text{new}}} \left(\frac{5}{6}\right)^{(d-o_{\text{new}})}$$

Optimality Equations: Probability of Player 2 Exceeding Player 1's Winning Score

Let $P_{\text{exceed}}(\Delta, o)$ denote the probability that player 2 will exceed player 1's score ≥ 50 where $\Delta = j - (i + k)$ (their score difference) and o is the number of rolled 1s set aside on player 2's final turn. Then,

$$P_{\text{exceed}}(\Delta, o) = \begin{cases} 0 & \text{if } o \geq 3 \\ 1 & \text{if } \Delta < 0 \\ \sum_{n=0}^{2-o} P_{\text{new1s}}(5 - o, n) P_{\text{exceed}}(\Delta - (5 - o'), o') & \text{otherwise} \\ & \text{where } o' = o + n \end{cases}$$

Optimality Equations: Probability of Winning with a Roll

The probability of winning with a roll $P_{\text{roll}}(p, i, j, k, o)$ under the assumption of optimal play thereafter is:

$$P_{\text{roll}}(p, i, j, k, o) = \begin{cases} P_{\text{exceed}}(j - i, o) & \text{if } p = 2 \\ & \text{and} \\ & j \geq 50 \\ \sum_{n=0}^{2-o} P_{\text{new1s}}(5 - o, n)P(p, i, j, k + 5 - o', o') + \\ \sum_{n=3-o}^{5-o} P_{\text{new1s}}(5 - o, n)(1 - P(3 - p, j, i, 0, 0)) & \text{otherwise} \end{cases}$$

A player can (and should) never hold at the beginning of the turn when the turn total is 0, so we express this by treating such rule-breaking as a loss. Thus, ...

Optimality Equations: Probability of Winning with a Hold, Roll/Hold Decision

the probability of winning with a hold $P_{\text{hold}}(p, i, j, k, o)$ under the assumption of optimal play thereafter is:

$$P_{\text{hold}}(p, i, j, k, o) = \begin{cases} 0 & \text{if } k = 0 \text{ or } (p = 2 \text{ and } j \geq 50, i) \\ 1 & \text{if } p = 2 \text{ and } i + k \geq 50, j \\ 1 - P(3 - p, j, i + k, 0, 0) & \text{otherwise} \end{cases}$$

Then the probability of winning $P(p, i, j, k, o)$ under the assumption of optimal play is:

$$P(p, i, j, k, o) = \max(P_{\text{roll}}(p, i, j, k, o), P_{\text{hold}}(p, i, j, k, o))$$

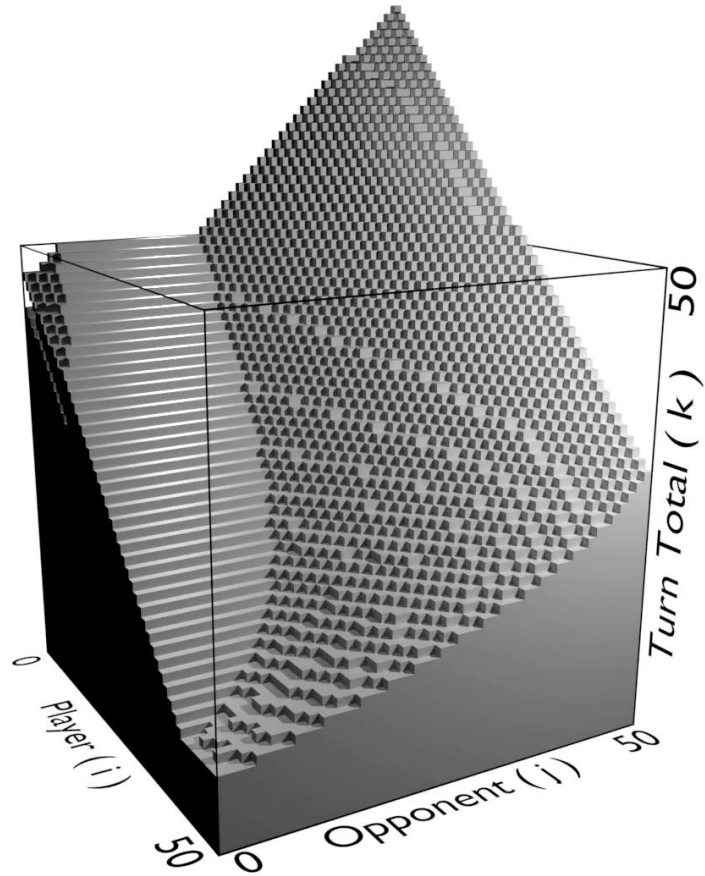
Solving Optimality Equations

- Equations (P_{new1s} , P_{exceed}) are solved through **dynamic programming** first.
- Cyclic, recursive P is solved through a **variation of value iteration**:
 - From initial arbitrary P estimates, substitute estimates in equation right-hand sides.
 - Compute the left-hand side P values as new, better estimates.
 - Terminate iterations of previous steps when the maximum change to a P estimate is $\leq 1 \times 10^{-14}$.

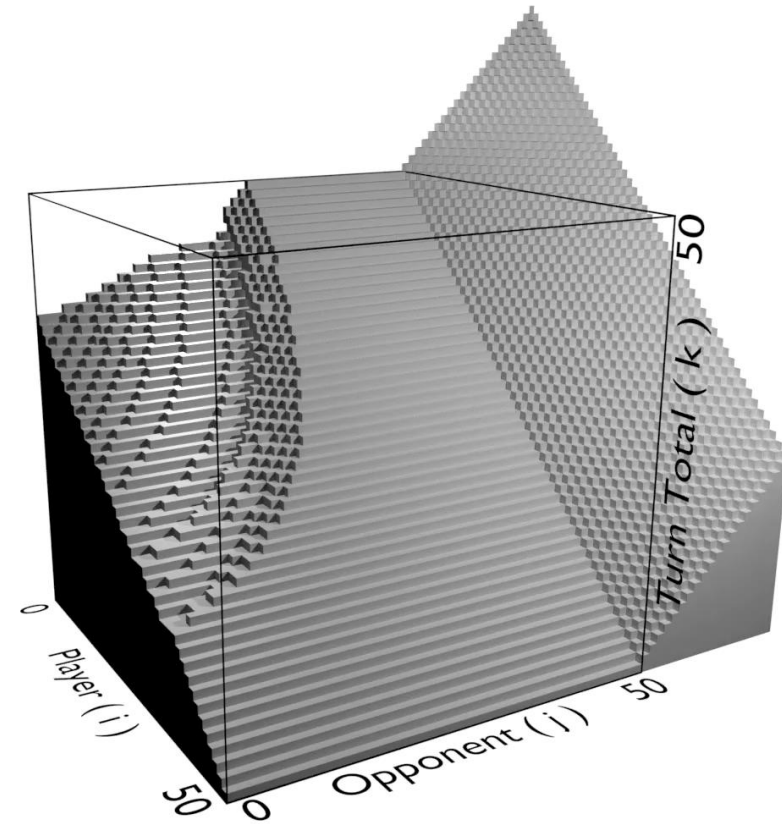
First Player Advantage and Compensation Points (Komi)

- Player 1 finishes with $\geq 50 \rightarrow$ Player 2 must exceed Player 1's score
- Player 2 has a knowledge advantage, knowing what score is needed to win.
- With optimal play, player 1 and player 2 have win rates of 0.4495 and 0.5505, respectively (a **10% gap!**).
- In the game of Go, “komi” are compensation points designed to make games more fair.
- In the Great Rolled Ones game, **player 1 should start with 3 compensation points (komi)**, bringing player 1's win rate up to 0.4955 (a **0.9% gap**) for most fair play.

Optimal Play (zero 1s set aside)

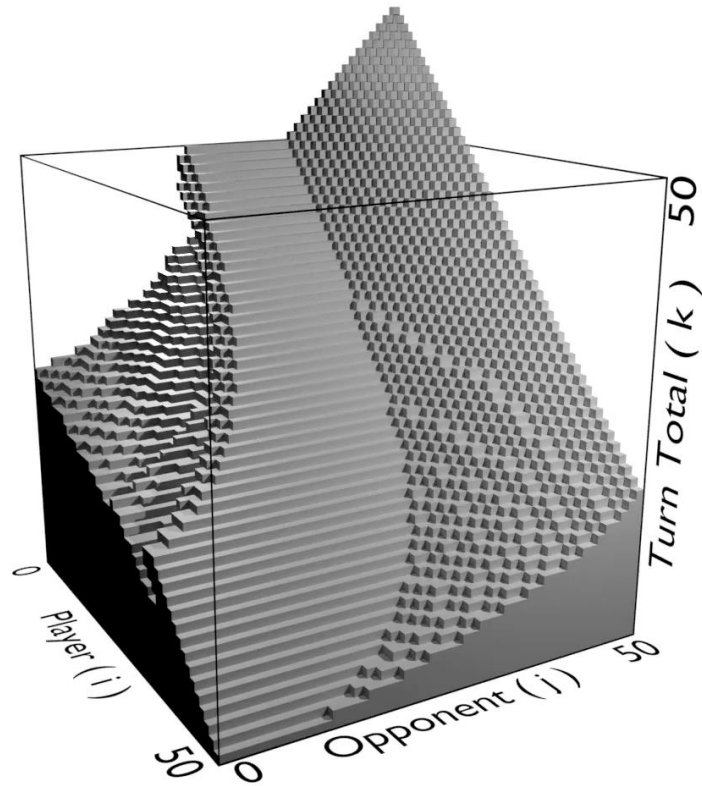


Player 1

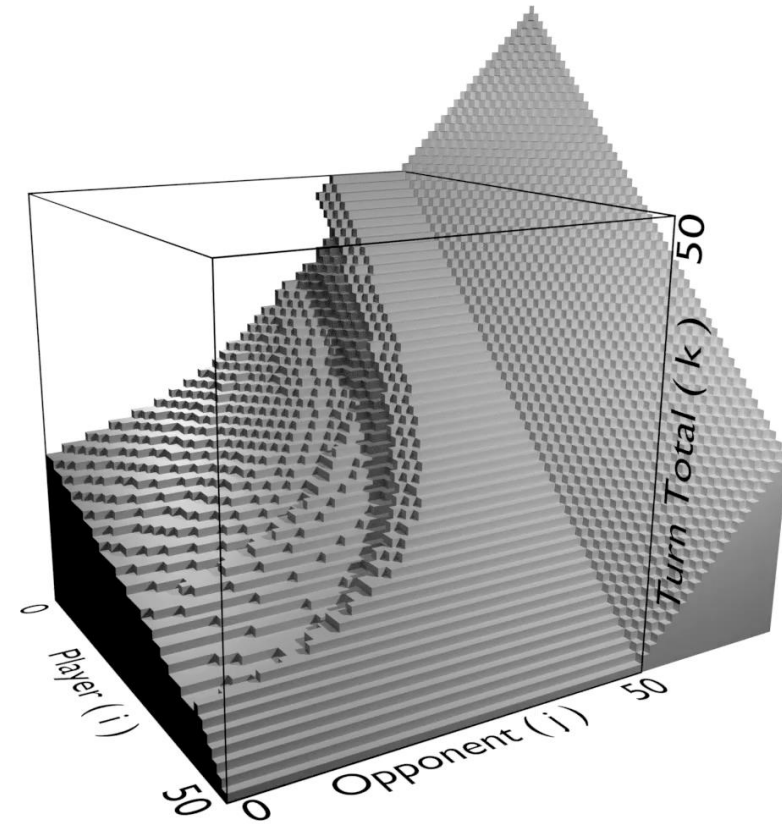


Player 2

Optimal Play (one 1 set aside)

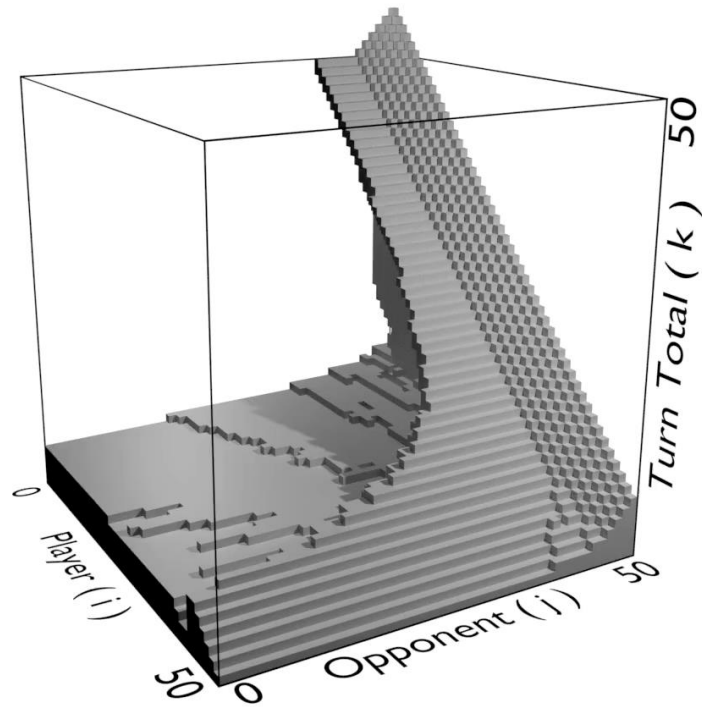


Player 1

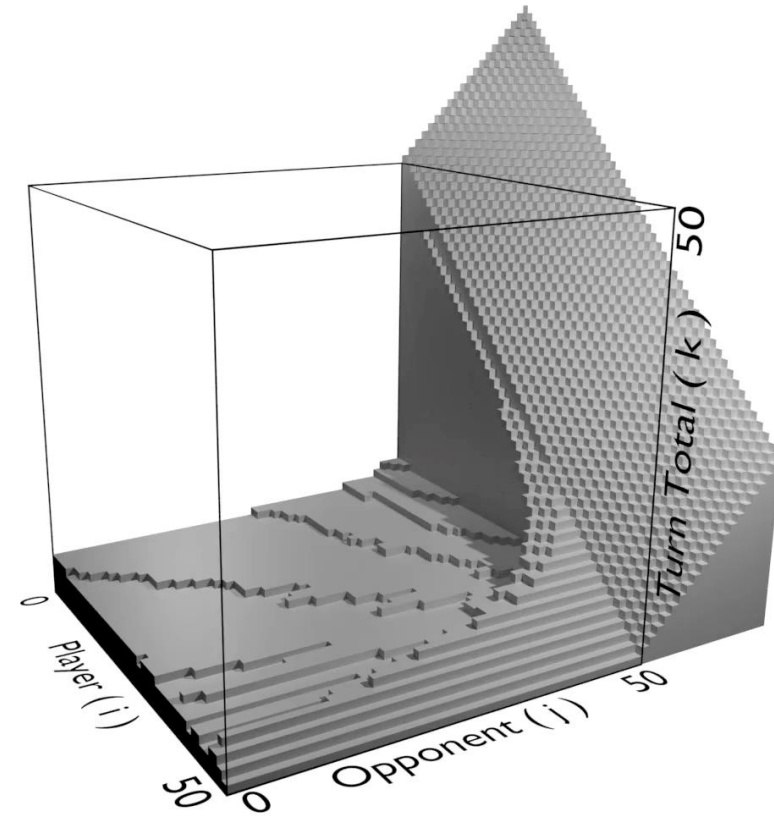


Player 2

Optimal Play (two 1s set aside)

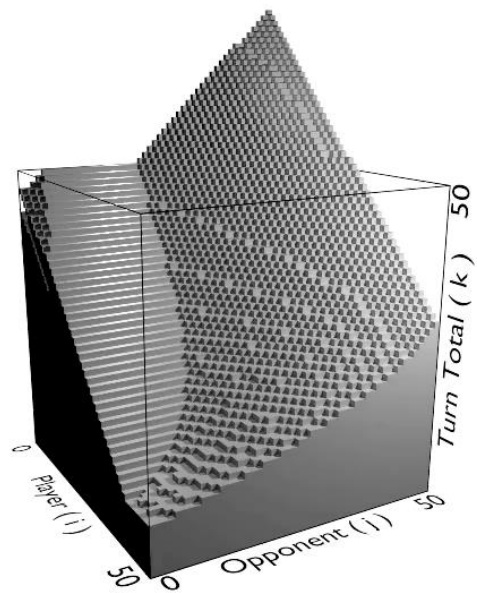


Player 1

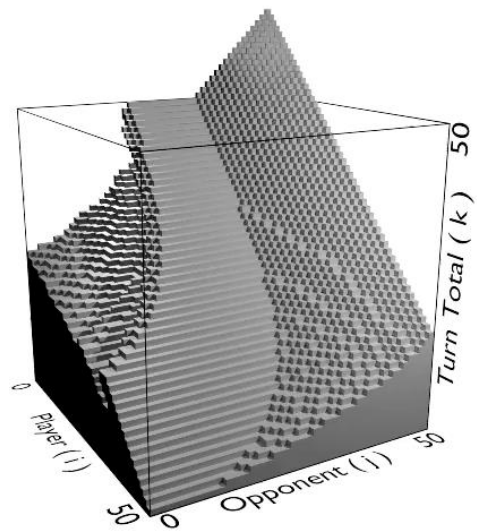


Player 2

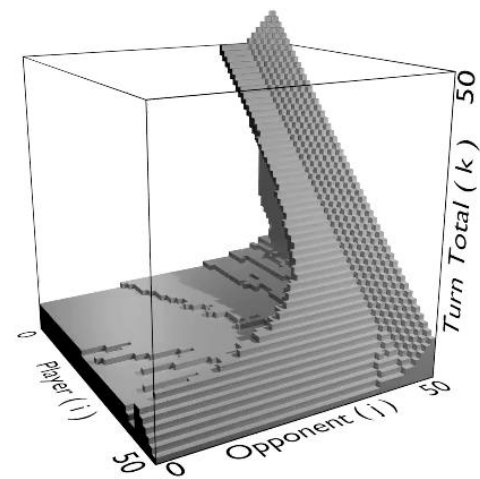
Zero 1s



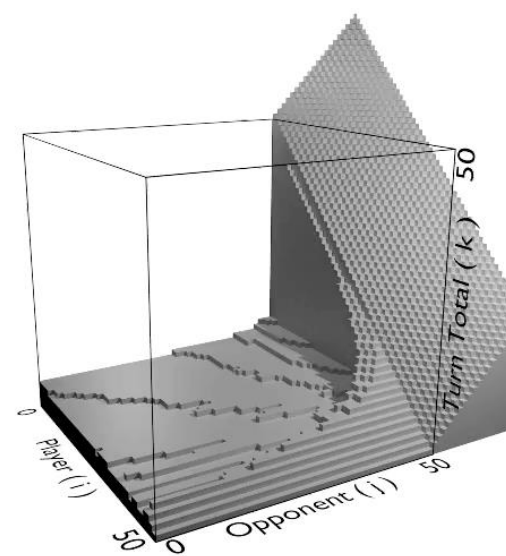
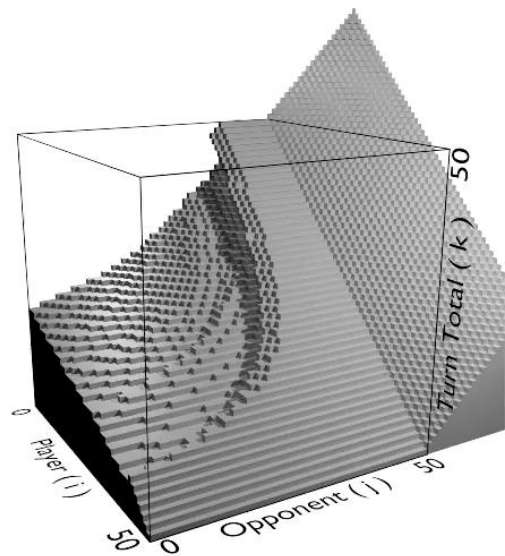
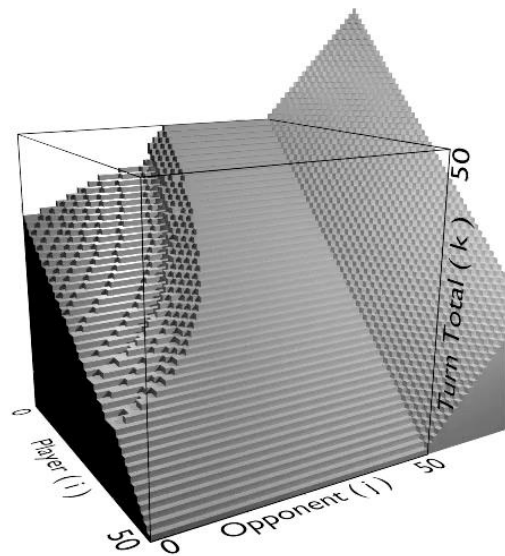
One 1



Two 1s



Player 1



Player 2

Human-Playable Policies

- By human-playable, we mean involving simple mental arithmetic and limited recall of cases and constants.
- We observe a continuum of tradeoffs from simple play policies with modest performance, to complex play policies with excellent performance.

Policy	Difference
Roll with 4 or 5 Dice	-0.0536
Fixed Hold-At	-0.0268
Simple Player and Ones Cases	-0.0201
Keep Pace, End Race, by Case	-0.0100

Fig. 2: Differences between human-playable and optimal policy win rates

Simple Player and Ones Cases (cont.)

- For player 1,
 - roll with 5 dice.
 - With 4 dice, hold at or beyond the goal with a lead of at least 20.
- For player 2,
 - if player 1 has reached the goal score, exceed it.
 - Otherwise, if player 2 can hold and win, do so.
 - Otherwise, player 2 always keeps rolling to win with 4 or 5 dice.
- With 3 dice, both players should hold if it reaches the goal score or if the turn total is at least 5.
- Such play wins only ~2.0% less than optimal play!

Conclusions

- Optimal play has been computed for the Great Rolled Ones game.
- 3 compensation points should be given initially to Player 1 for greatest fairness
- Among the variety of human playable strategies analyzed, we shared the “Simple Player and Ones Cases” strategy that has a 2% gap from the optimal win rate:
 - (Player 2 must exceed a winning Player 1 score.)

1s Rolled	Player 1	Player2
0	Roll and do not hold	Roll to win
1	Hold at ≥ 50 points with a lead of ≥ 20	Roll to win
2	Hold at ≥ 5 turn total or ≥ 50 points	Hold at ≥ 5 turn total or ≥ 50 points

Roll with 4 or 5 Dice

Algorithm 1: Roll with 4 or 5 dice

Input : player p , player score i , opponent score j , turn total k , ones rolled o

Output: whether or not to roll

```
1 if  $p = 2 \wedge j \geq 50 \wedge i + k \leq j$  then           // player 2 must exceed player 1
2 |   return true
3 else if  $p = 2 \wedge i + k \geq 50$  then           // player 2 must hold at goal score
4 |   return false
5 else                                           // roll with 4 or 5 dice
6 |   return  $o < 2$ 
7 end if
```

Fixed Hold-At

Algorithm 2: Fixed hold-at

Input : player p , player score i , opponent score j , turn total k , ones rolled o

Output: whether or not to roll

```
1 if  $p = 2 \wedge j \geq 50 \wedge i + k \leq j$  then           // player 2 must exceed player 1
2   |   return true
3 else if  $i + k \geq 50$  then                           // player 2 holds and wins
4   |   return false
5 else if  $o = 0$  then                                   // keep rolling with 5 dice
6   |   return true
7 else if  $o = 1$  then                                   // hold at 24 with 4 dice
8   |   return  $k < 24$ 
9 else                                                  // hold at 4 with 3 dice
10  |   return  $k < 4$ 
11 end if
```

Simple Player and Ones Cases

Algorithm 3: Simple player and ones cases

Input : player p , player score i , opponent score j , turn total k , ones rolled o

Output: whether or not to roll

```
1 if  $p = 1$  then                                     // player 1 cases
2   | if  $o = 0$  then                                   // keep rolling with 5 dice
3   |   | return true
4   | else if  $o = 1$  then                             // hold at goal with  $\geq 20$  lead with 4 dice
5   |   | return  $k < \max(50 - i, 20 + j - i)$ 
6   | else                                           // hold at 5 or goal with 3 dice
7   |   | return  $k < \min(50 - i, 5)$ 
8   | end if
9 else                                             // player 2 cases
10  | if  $j \geq 50$  then                               // player 2 must exceed player 1
11  |   | return  $i + k \leq j$ 
12  | else if  $i + k \geq 50$  then                     // hold at goal score
13  |   | return false
14  | else if  $o < 2$  then                             // roll with 4 or 5 dice
15  |   | return true
16  | else                                           // hold at 5 or goal with 3 dice
17  |   | return  $k < \min(50 - i, 5)$ 
18  | end if
19 end if
```

2.0% gap

Keep Pace, End Race, by Case

Algorithm 4: Keep pace, end race, by case

Input : player p , player score i , opponent score j , turn total k , ones rolled o
Output: whether or not to roll

```
1  $\delta \leftarrow j - i$ 
2 if  $p = 1$  then // player 1 cases
3   if  $o = 0$  then // hold at goal with  $\geq 38$  lead with 5 dice
4     return  $k < \max(50 - i, 38 + \delta)$ 
5   else if  $o = 1$  then
6      $h \leftarrow 22 + \delta$  // hold with a  $\geq 22$  lead with 4 dice
7     if  $i \geq 10 \vee j \geq 23$  then
8       // if player 1 / 2 has scored 10 / 23, resp.
9        $h \leftarrow \max(50 - i, h)$  // then at least roll for the goal
10    end if
11    return  $k < h$ 
12  else if  $i + j \geq 71$  then
13    // reach the goal when the player score sum reaches 71
14    return  $k < 50 - i$ 
15  else // hold at 5 or goal with 3 dice
16    return  $k < \min(50 - i, 5)$ 
17  end if
18 else // player 2 cases
19   if  $j \geq 50$  then // player 2 must exceed player 1
20     return  $k \leq \delta$ 
21   else if  $o = 0$  then // keep rolling with 5 dice
22     return true
23   else if  $o = 1$  then // with 4 dice
24     if  $i \geq 20 \vee j \geq 32$  then
25       // if player 1 / 2 has scored 20 / 32, resp.
26       return  $k < 50 - i$  // then roll for the goal
27     else // else hold with  $\geq 28$  lead
28       return  $k < 18 + \delta$ 
29     end if
30   else if  $i + j \geq 84$  then
31     // reach the goal when the player score sum reaches 84
32     return  $k < 50 - i$ 
33   else // hold at 5 or goal with 3 dice
34     return  $k < \min(50 - i, 5)$ 
35   end if
36 end if
```

1.0% gap

Future Work

- Supervised learning of win probabilities for nonterminal states could compress our precise tabular computation.
- One-step backup of approximate win probabilities would likely yield excellent roll/hold decisions.
- Question: How well would different models/techniques perform for trading off performance for reduced memory requirements?